COMP2310/COMP6310 Systems, Networks, & Concurrency

Convener: Shoaib Akram



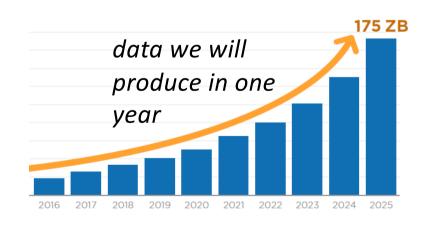
Shoaib Akram

School of Computing, (Jan 2020 –)

Ph.D., 2019

Teaching: Computer Organization & Program Execution (semester 1) **Systems, Network, and Concurrency** (semester 2)

Research: Fast analytics over big datasets with a focus on emerging memories and fast storage devices







Quick Logistics

Course webpage: https://comp.anu.edu.au/courses/comp2310/

Lectures (on the website)

- Lecture slides
- Lecture videos (*Echo360*)
- 2 hours reserved (some lectures may be shorter, demos etc)

Policies (will be up shortly)

General conduct, assignment submissions, support, management, grading

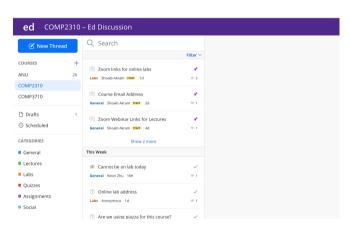
Resources

- Past exam with solution and rubric (up already)
- Stuff needed to finish the labs and assignments

Edstem

We will use edstem for all communication

- If you ignore edstem, you will miss key announcements
 - Drop-in sessions, make-up lectures, problems, exercises, corrections, lecture timing
 - Ask questions on edstem first (most likely you will receive a response quickly)
- Ask instructors private questions on edstem
- Students are added/dropped automatically



Course Email

comp2310@anu.edu.au

- Do not send me a direct email except for requests:
 - Super urgent
 - Personal
 - EAP-related

Motivation

Recall: How do we make electrons do the work?

Problem Statement: "Save the planet"
The Algorithm
Program in a High-Level Language
Instruction Set Architecture (ISA)
Microarchitecture
Circuits
Devices ~

Recall: How do we make electrons do the work?

- Using a sequence of systematic transformations
 - Developed over six decades
- Each step must be studied and improved for the whole stack to work efficiently

Recall: Transformation Hierarchy

- We call the steps of the process: Levels of transformation OR Transformation hierarchy
- At each level of the stack, we have choices
 - Language: Java, Python, Ruby, Scala, C++, C#
 - ISA: ARM, x86, SPARC, PowerPC, RISC-V
 - Microarchitecture: Intel, AMD, IBM
- If we ignore any of the steps, then we cannot
 - Make the best use of computer systems
 - Build the **best** system for a set of programs

Problem

Algorithm

Program

Architecture

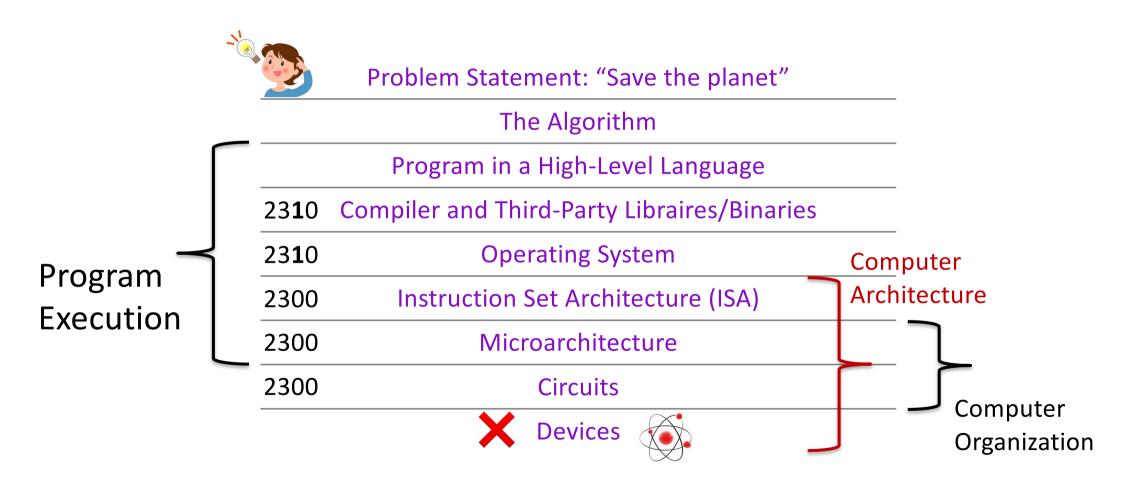
micro-arch

circuits

devices



Recall: Transformation Hierarchy & Us



Recall: Hardware and Software

Proble

Problem Statement: "Save the planet"

Software

The Algorithm

Program in a High-Level Language

Compiler and Third-Party Libraires/Binaries

Operating System

ISA = Hw/Sw

boundary/interface

Instruction Set Architecture (ISA)

Microarchitecture

Circuits

Devices

Recall: Two Recurring Themes

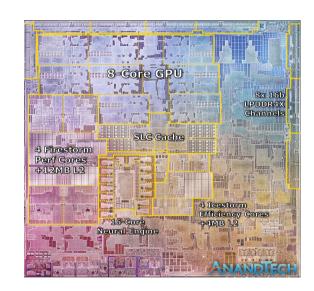
The notion of abstraction

Hardware versus software

Recall: The Notion of Abstraction

Abstraction: Know components from a high level of detail

No human (programmer) can track 10 billion elements. Computer systems work because of abstraction!



Apple M1 Chip
Billions of transistors
All working in parallel

Recall: The Notion of Abstraction

- Abstraction: View the world from a higher level
- Focus on the important aspects
 - Input? Output? X = ADD or MULTIPLY



- Raise the level of abstraction for productivity and efficiency
- But what if the world below does not work as expected?
 - To deal with it, we need to go below the abstraction layer
- Deconstruction: To un-abstract when needed
 - Important skill

Recall: The Notion of Abstraction

- We will use this theme a lot!
 - Each layer in the transformation hierarchy is an abstraction layer!



Algorithm

Program

Architecture

micro-arch

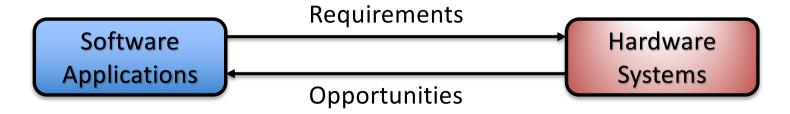
circuits

devices



Recall: Hardware versus Software

- Hardware versus software
 - Hardware: Physical computer
 - Software: Programs, operating systems, compilers
- One view: Ok to be an expert at one of these
- Hw and Sw: Two parts of the computer system
 - COMP2300 view: Knowing the capabilities/limitations of each leads to better overall systems

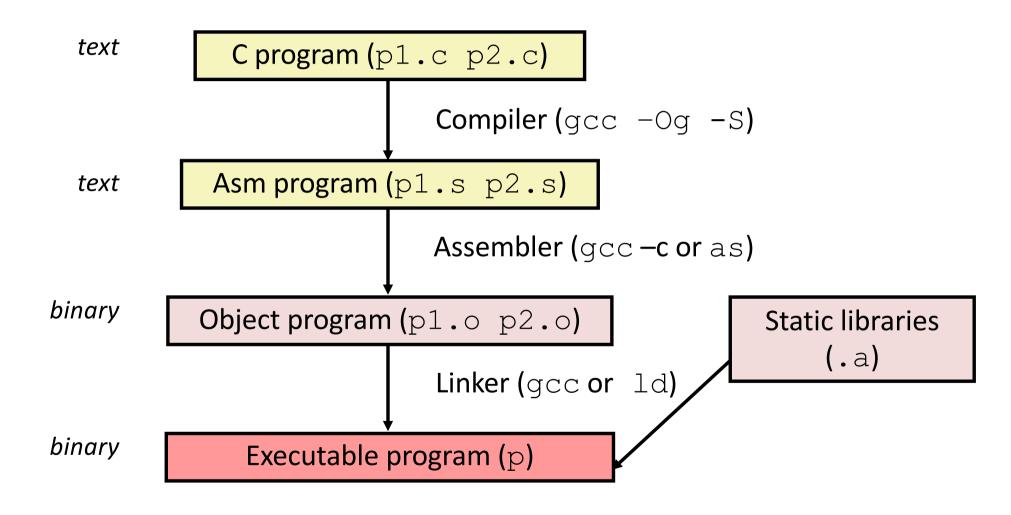


COMP2310 deepens this knowledge

Role of Compiler

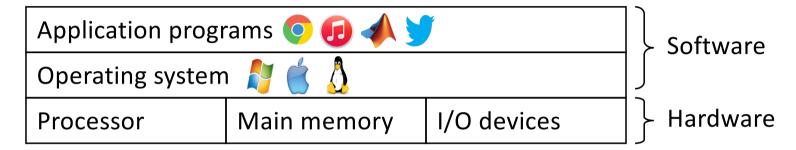
- What does a compiler do?
 - Translates high-level code into assembly
- More generally, the compilation toolchain generates machine code in a sequence of stages:
 - translate a group of related source files into assembly
 - resolve inter-dependencies between source files (linking)
 - handle the *linking* of any external libraries
 - perform optimizations (make use of special hardware features)
- It is more complex than line-by-line C to assembly translation
- Learning the process is important from a performance, efficiency, security, and hacker perspective

Turning C into Object Code (details later)



Role of Operating System (OS)

Operating system



- Enables safe abstractions of hardware resources
- Virtualizes hardware for use by programs
 - Gives each program the illusion that it has the entire resource for itself
- Manages the hardware resources for efficient and safe working of the system

COMP2310, Goal # 1

- Deepen the understanding of how applications interact with compiler and OS, and hardware
- Today, critical for software to be correct, performant, efficient, secure
- Demystify how programs are loaded into memory and executed
 - What happens when you click an icon to start an application?
 - Or type the program name into a shell program and press enter

COMP2310, Goal # 1 (cont'd)

- How are C programs translated into x86-64 assembly?
 - Compilation and linking fundamentals
 - Object files, executable formats, etc
 - Implementation of loops, procedure calls, data structures (reprise)
 - Optimizations done (not done) by the compiler

Assembly is Important!

Intel x86-64 ISA widely used in server hardware



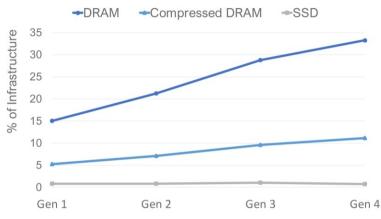
- Tuning performance
 - Understanding optimization done (not done) by the compiler
 - Understanding behavior of programs and exploiting choice via compile-time options
- Writing systems software (device drivers)
- Fighting security vulnerabilities
- Behavior of buggy programs

COMP2310, Goal # 1 (cont'd)

- What does the memory hierarchy of looks like?
 - How do caches work in more detail?
 - What is their impact on program behavior? (programmer's perspective)
 - How does main memory differ from a disk drive?
 - How does device behavior impact the design of computer programs?

Memory matters!

- Memory is a limited resource
 - Must be carefully managed
- Memory bugs are hard to detect
 - Understanding pointers and memory allocators helps
- Memory performance is not always uniform
 - Caches, virtual memory effects need to understood



Meta datacenters, 2022

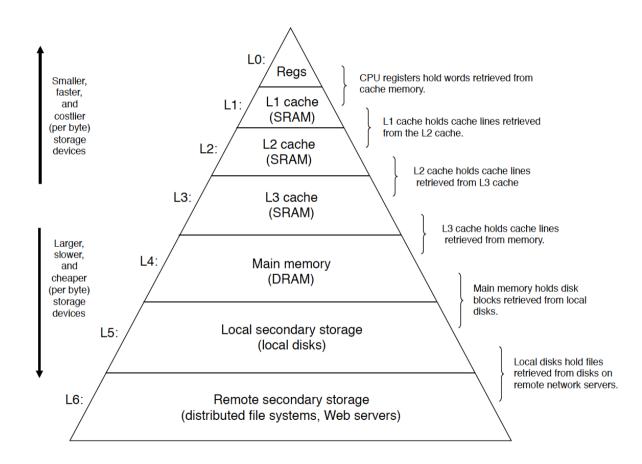
Memory matters!

81.8ms

4.3ms 2.0 GHz Intel Core i7 Haswell

- Hierarchical memory organization
- Performance depends on access patterns
 - Including how step through multi-dimensional array

Memory Hierarchy



COMP2310, Goal # 1 (cont'd)

- How does the operating system abstract hardware resources for use by application programs?
 - Processes
 - Virtual memory
 - Files
- All these are abstractions the OS uses to isolate computer programs from each other
 - Our focus is NOT on (re)building these "mechanisms" but write programs to use them

COMP2310, Goal # 2

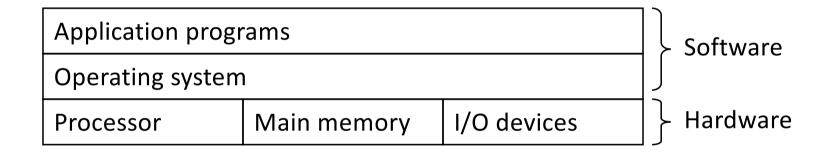
- Understand how applications use the operating system (OS) and the C standard library for writing real-world applications
- Search engines
- Databases
- Android memory manager

COMP2310, Goal # 2

- Write low-level code that interfaces with the operating system kernel and C library
- What interface (and interesting system calls) does the operating system provide?
- Learn to
 - Implement memory allocators, read/write from/to storage disks and SSDs
 - Communicate with the outside world (networking), and manage concurrently running processes and applications

User Code vs. Kernel Code

 OS manages the hardware, interposed b/w the program and hardware



- Application code that runs on top of OS or any resource manager in general is user code (or user program)
- The code that manages the hardware is kernel code
- CPU is either in user mode or kernel mode

What is OS kernel?

- Core component of OS that manages the hardware
 - Device management (keyboard, mouse, display, etc)
 - Memory (RAM) management
 - Network management
 - Storage management
 - Filesystem code
- What else is in the OS?
 - Shell
 - GUI
 - Utilities

What does an OS do?

- Manages the hardware, interposed b/w the program and hardware
- Our high-level view of system (no network for simplicity)



OS manages CPU (processor), memory (RAM), input/output (I/O) devices (keyboard, disk, display, network), and files on disk

How do applications use the OS?

- OS provides services to be accessed by user programs
- Programs can make use of "system calls" on Linux and Windows application programming interface ("API")
 - Allocate memory for me
 - Read "N" bytes from file F into memory location "M"
 - Write "N" bytes from memory location "M" into file F
 - Establish a network connection to <u>www.anu.edu.au</u>
 - Write "N" bytes to the network connection
 - Put me to sleep

How applications use the OS?

- OS provides an interface for applications to use
 - Programs access hardware/device capabilities through this interface
 - Different hardware → Same interface
 - Interface is constant, its implementation is OS specific
- We need to learn this interface to write interesting applications
 - Learning "just enough" details of the implementation to write correct, efficient, secure programs

Goal # 3, Systems Programming

- All these aspects will help you become a systems programmer
- Systems programmers
 - write low-level tools such as compilers, operating systems, and debuggers
 - they must have an acute awareness of the environment, e.g.,
 Linux versus Windows
 - they must use system calls for the specific OS
 - contrast with Python, Ruby, Java programs for business or ML
 - high-level libraries abstract OS and hardware details
 - C library abstracts OS/hardware but many Linux C programs interface with the kernel API

Example of Pure User-Level C Program

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#define LEN 10000000
// struct of arrays
// all i points are stored in a single contiguous array
// all i points are stored in a single contiguous array
// all k points are stored in a single contiguous array
struct pointarray3D {
        int i[LEN];
        int i[LEN]:
        int k[LEN];
};
struct pointarray3D points;
int sum k (struct pointarray3D *points) {
        int sum = 0;
        return sum:
int main() {
        for (int idx = 0; idx < LEN; idx++) {
                points.i[idx] = 1;
                points.j[idx] = 1;
                points.k[idx] = 1;
        int sum = sum_k(&points);
        printf("sum of all k points is %i \n", sum);
        return 0:
```

- Will not crash your machine if you did something wrong
- Programmer's creativity is more critical in solving the problem
 - Can get by not knowing how array looks like in memory
- Uses a C library function for printing to the screen
- C library takes care of making it happen for the programmer

Example of Pure System-Level C Program

```
#include de linux/module.h>
#include linux/moduleparam.h>
#include inux/init.h>
#include linux/kernel.h>
#include linux/proc fs.h>
#include <asm/uaccess.h>
                           "abcdefghijklmnopqrstyvwxyz"
#define alpha
#define BUF_SIZE
              Do while having super-user privillages:
                   insmod mine.ko'
                                            insert module into the kernel.
                 ' rmmod mine.ko' :
                                            remove module from the kernel.
// The entry will be created into the ' /proc ' directory // the directory that will hold the new device.
static struct proc_dir_entry *ent;
static char message[BUF SIZE];
static ssize_t mwrite(struct file *file, const char __user *ubuf, size_t count, loff_t *offset)
         unsigned int i;
         int rv;
         char user *p = ubuf;
        printk(KERN_INFO "Write Handler\n");
printk(KERN_INFO "Size: %d , Offset: %d \n", count, *offset);
        if(count > BUF SIZE)
                 return -EFAULT;
         rv = copy_from_user(message, p, count);
         printk(KERN_INFO "Byte not copied: %d\n",rv);
         printk(KERN INFO "device have been written\n");
static ssize t mread(struct file *file, char user *ubuf, size t count, loff t *ppos)
         char __user *ptr;
         printk(KERN INFO "Read Handler\n");
         printk(KERN ALERT "Count : %d\n",count);
         ptr = ubuf:
         if(count > BUF SIZE) {
                  printk(KERN_INFO "Adjusting size:\n");
                  count = BUF SIZE;
                  printk(KERN INFO "Size: %d\n",count);
                  return (-1);
         copy to user(ubuf, message, count);
        return count:
// File operations.
static struct file_operations fops =
         .owner = THIS MODULE,
         .read = mread.
         .write = mwrite.
```

- Device driver code
- Most likely crash your machine if you did something wrong
- Requires intricate knowledge of the hardware for which driver is being written
- Uses Linux kernel sources to reuse functionality
- Even "printf()" is not available
- No C library

Example of User-Space System-Level C Program

Focus of this course!

```
1 #include <unistd.h>
 2 #include <string.h>
   char ptype[10];
8 int main()
        int size = 50 * sizeof(int);
       void *addr = mmap(∅, size, PROT READ | PROT WRITE, MAP SHARED | MAP ANONYMOUS, -1, ∅);
       printf("Mapped at : %p \n\n", addr);
       int *shared = addr;
       pid t fork return = fork();
       if (fork_return > 0)
           shared[0] = 40;
           shared[1] = -20;
                  (ptype, "Parent");
           int status;
           waitpid(-1, &status, 0);
            printf("Child : shared[0] = %d , shared[1] = %d \n", shared[0], shared[1]);
           shared[1] = 120;
            strcpy(ptype, "Child ");
       printf("%s : shared[0] : %d\n", ptype, shared[0]);
        printf("%s : shared[1] : %d\n", ptype, shared[1]);
       munmap(addr, size);
        return 0;
35 }
```

- Won't crash your machine
 - But program is likely to crash if something is wrong
- Uses system call wrappers provided by C library
- Uses "interesting" system calls
 - fork() spawn a new virtual CPU
 - mmap() instantiates a region in the process' address space

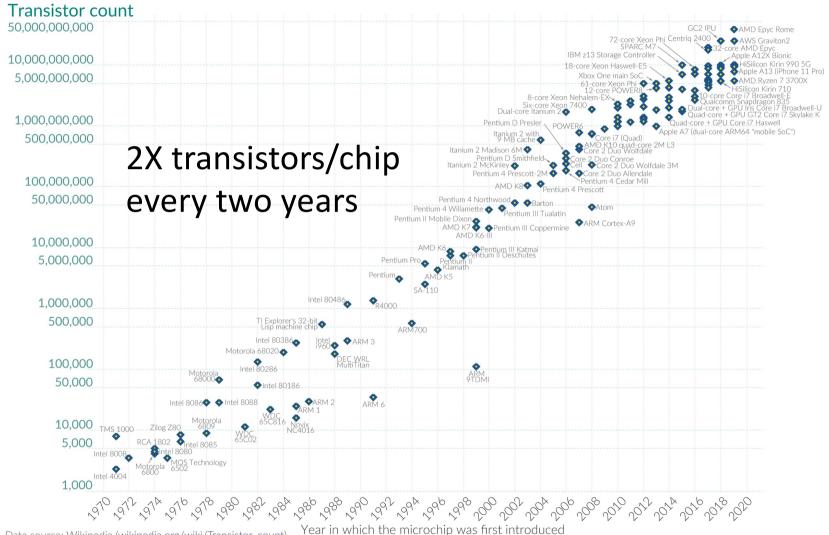
CPU Trends



Moore's Law: The number of transistors on microchips doubles every two years



Moore's law describes the empirical regularity that the number of transistors on integrated circuits doubles approximately every two years. This advancement is important for other aspects of technological progress in computing – such as processing speed or the price of computers.



Data source: Wikipedia (wikipedia.org/wiki/Transistor count)

OurWorldinData.org - Research and data to make progress against the world's largest problems.

Licensed under CC-BY by the authors Hannah Ritchie and Max Roser.

End of Dennard Scaling

 Dennard scaling: As transistors get smaller, their density stays constant



- In every technology generation, the area and power consumption of individual transistors is halved
 - With twice the number of transistors, power consumption still stays the same

Dennard scaling broke down b/w 2005-2007

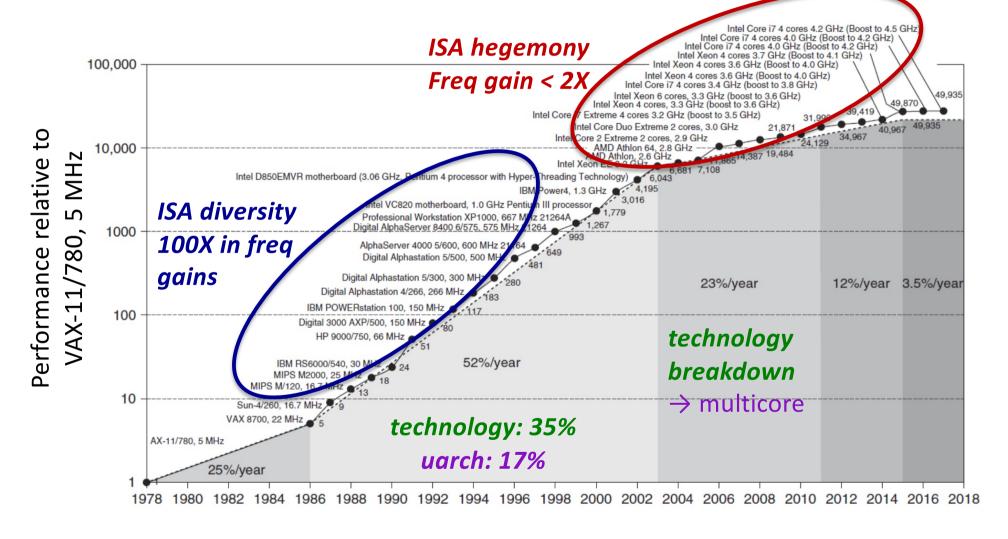
→ As we add more transistors, power consumption for for a chip with the same area increases

End of Dennard Scaling

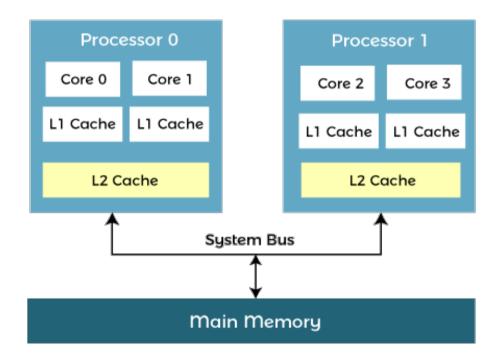
Implication: Frequency cannot increase any further because that would make the power problem even worst → Industry shifted to multicores!

```
https://silvanogai.github.io/posts/dennard/
https://www.maketecheasier.com/why-cpu-
clock-speed-isnt-increasing/
```

Uni-Processor Performance



Modern System



Software must exploit parallelism for performance

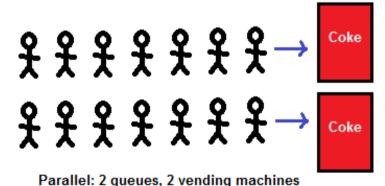
Concurrency and Parallelism

- Concurrency: When the execution of two programs overlap in time
 - Concurrency has always been important
 - Multiple users time-sharing a uniprocessor system
 - Process an incoming request from the network,
 while the user is watching a video recording
- Parallelism: When two programs use dedicated resources (e.g., two separate CPUs) to run at the same time
 - Multicores have made parallelism critical to get more performance out of modern hardware

Concurrency and Parallelism Real-Life Example



Concurrent: 2 queues, 1 vending machine



Managing System Resources

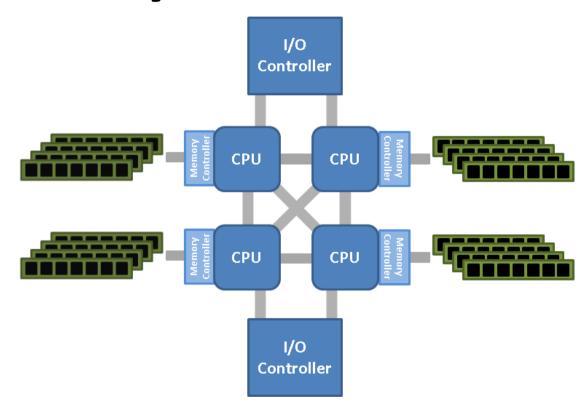
- Many interesting debates in computer systems
 - RISC vs. CISC
 - Compiler vs. hardware exploitation of ILP
 - Manual (C++) vs. automatic memory management (Java)
- One more debate
 - Should "certain" hardware features be exposed to user-level application or not?
 - One camp: User-level programmer possesses better knowledge of application logic than hardware or compiler or OS
 - They can "tune" the feature to make the optimal use of it
 - Other camp: They may also do something wrong
 - Leave it to the hardware or compiler or OS

Managing System Resources: Examples

- CPU registers are exposed to software (OS and user-level)
- CPU caches are managed by hardware
 - We say caches are transparent to software
- A feature X is exposed to software, but OS utilizes the feature and user-level code has no way to access it
 - Feature X is transparent to user-level code
 - Feature X is visible to OS, or X is exposed to OS
 - Physical memory is an example (what? that's why 2310 exists!)

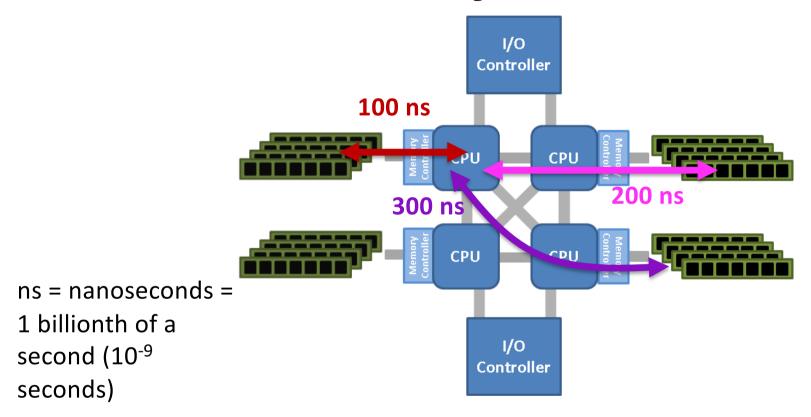
Modern NUMA System

Some memories are closer to the CPU, while others are far away. Wrong data placement can hurt performance.



User-space software must be aware of Non-Uniform Memory Access (NUMA) architectures (one view)

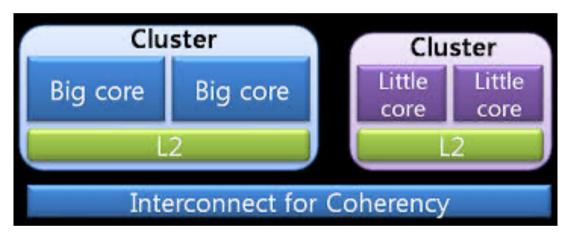
Modern NUMA System



User-space software must be aware of Non-Uniform Memory Access (NUMA) architectures (one view)

Another system (cell phones)

ARM big.LITTLE



Power-hungry Performance-driven

mission-critical tasks

Energy-efficient

background tasks

Why learn systems programming?

- Key takeaway: As a systems programmer, you can advise the OS (or any resource manager) to make the best use of the underlying hardware
- We will teach you how you can build applications that hook up with the kernel and do just that and other interesting things
- We won't build: CPU, OS, compiler, here

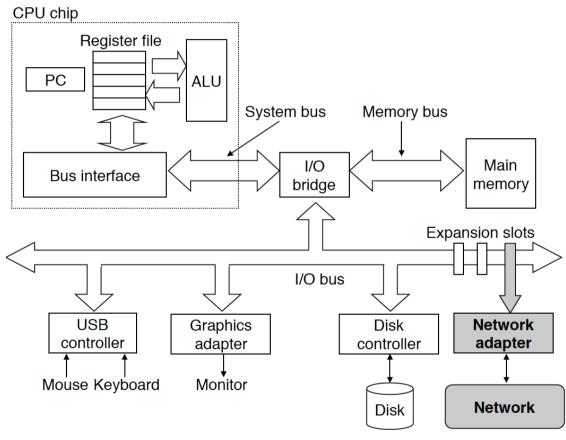
More Examples of System Features

- NUMA
- Heterogeneous multicore processors (e.g., big.LITTLE)
- Persistent memory (e.g., Intel Optane Persistent memory)
- CPU-FPGA platforms
- Computational storage devices (CSD)
- Programmable network interface cards
- Hyperthreading
- Turbo boosting, low-power modes
- Single Instruction Multiple Data (SIMD) instructions
- ML accelerators
- Some recent additions to ISA for hardware cache mgmt.
- Dynamic voltage and frequency scaling (DVFS)
- Processing in Memory (PIM)
- Heterogeneous-ISA multicore processors
- Remote memory
- Single-ISA multicore processors
- Intel Cache Allocation and Monitoring Technology
- Memory-semantic solid-state drives (SSDs)
- CXL-based memory expansion
- Software defined storage

Networking

- Web, social media, email, online games, all use the networking
- We will learn the basics of client-server model
- Writing simple networking applications in C

COMP2310: Holistic View of Computer System



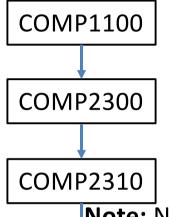
Course Perspective

- Most systems courses are Builder-Centric
 - Computer Organization (COMP2300), Microarchitecture
 - Build a CPU. Implement an ISA
 - Operating Systems (COMP3300, Alwen Tiu)
 - Implement portions of operating system
 - Compilers (COMP3710, Tony Hosking)
 - Write compiler for a simple language
 - Computer Networks (COMP3310)
 - Implement and simulate network protocols

Course Perspective

- COMP2310 is programmer-centric
 - By knowing more about the underlying system, you can be more effective as a programmer
 - Enable you to
 - Write programs that are more reliable and efficient
 - Incorporate features that require hooks into OS
 - E.g., concurrency, signal handlers
 - Things you will not see elsewhere or are required background knowledge
 - Not a course for **dedicated** hackers
 - We aim to bring the hidden hacker inside you!

Role within CS Curriculum



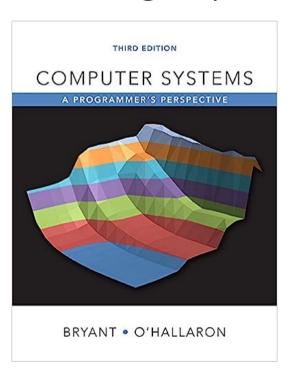
Note: Not a pre-req for all courses by ANU policy

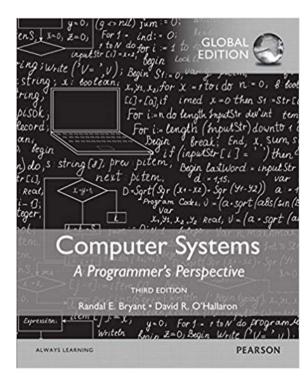
- Software Security
- Operating Systems
- Compilers
- Computer Networks
- High Performance and Scientific Computing
- Parallel Systems
- Computer Graphics
- Algorithms
- Databases

Content & Topics

Primary Textbook

- Textbook really matters for the course (problems, lectures, labs)
 - Textbook is not "just" a recommendation
- Warning: Paperback international version has "some" errors



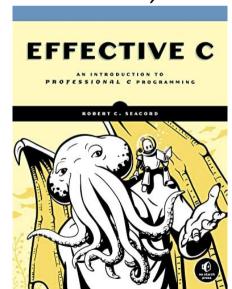


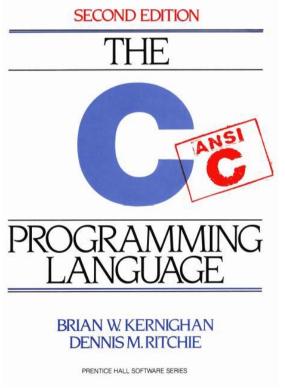
Useful Books on C (optional)

Kernighan & Ritchie, The C Programming Language, 2nd Edition

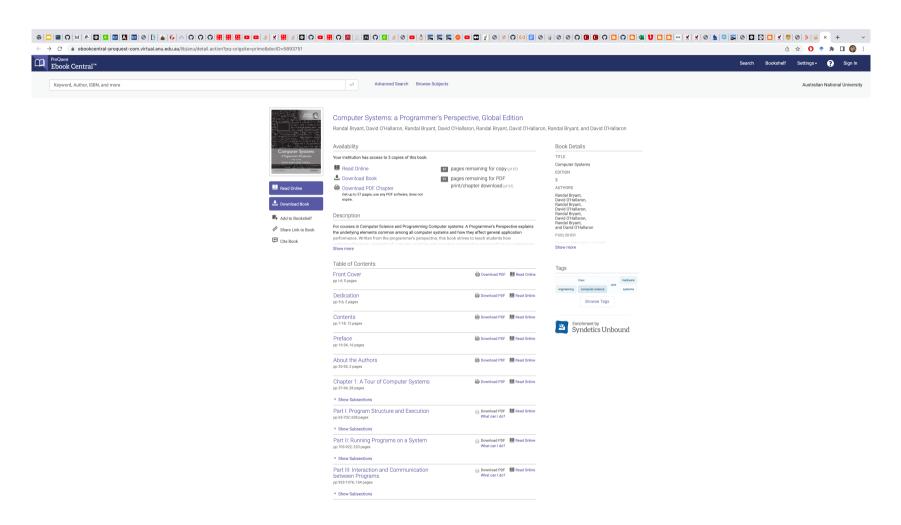
- "ANSI" (old-school) C
- Not too serious about things we now consider critical

Slightly advanced, more practical advice, modern





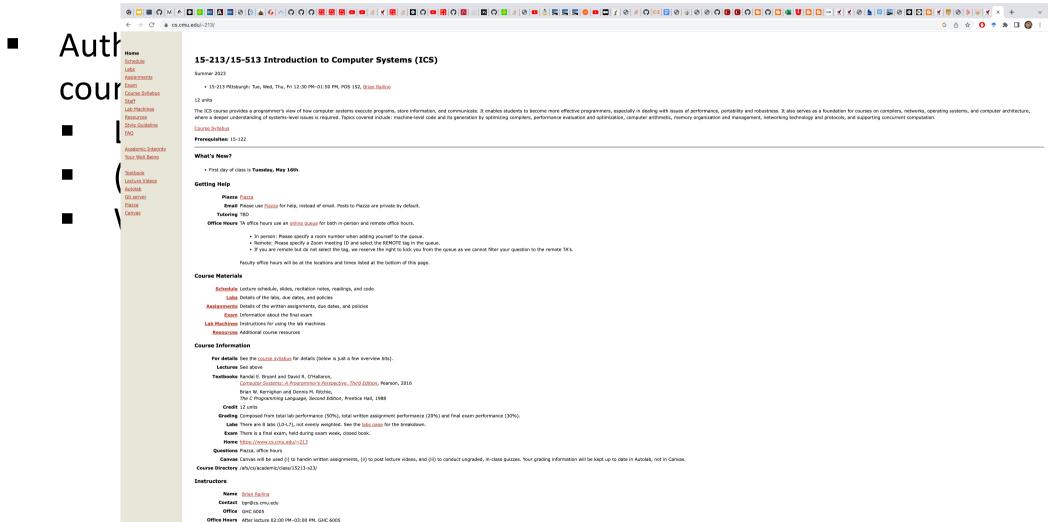
Textbook: Electronic edition available for ANU students



CMU 213

- Authors of the book at the Carnegie Mellon University created a course to accompany with the book
 - Lecture slides, problem sets, exams, labs, etc
 - (Acknowledgement) We use the material from the course
- We encourage you to explore the CMU course website
 - Note: Their course combines aspects of COMP2300 and COMP2310 into one course
 - Their starting point: COMP2300 starting point
 - Their CPU coverage is limited (programmer's perspective)
 - Key Point: Do not ignore COMP2310 & blindly follow CMU213

CMU 213



High-Level to Low-Level Translation

- C programming to x86-64 assembly
- Compilation steps
- Array allocation and access
- Heterogenous data structures
- Optimizations
- Security vulnerabilities

COMP2310 is not a C Programming Course

- Emphasis is on program transformation
- How does high-level code look in assembly?
- Do compilers always do the right thing?
- Programmers WILL write more efficient code if they have insight into transformation steps
- The power to reverse engineer object code and binaries
 - A.k.a. hackers! Security professionals' bread and butter

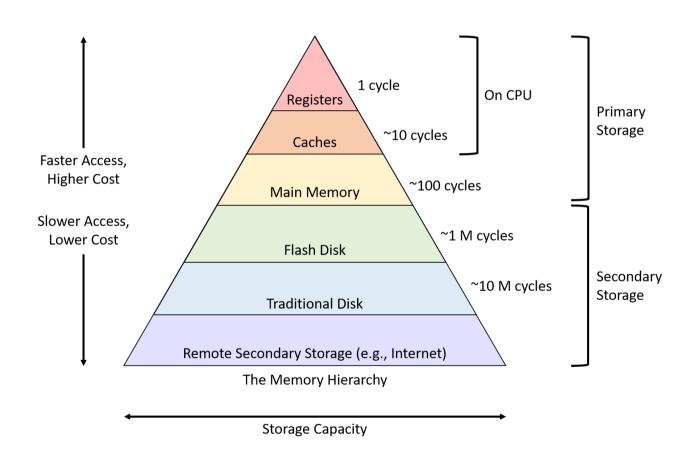
Qualified Answers to C Questions

- What are pros and cons of programming in C?
- Why should you NEVER use C in 2022? Why should everyone learn C (and then program in whatever language they like?
- Why is C insecure and what can be done about that?
- Why is Linux OS written in C? And many other datacenter software stacks?
- Why is C dominant in the embedded domain?

Exceptional Control Flow

- Processes
 - The illusion that each program has the entire CPU for its own use even though many programs might be co-running
- Exceptions and signals
- Address spaces
- How does Unix-like systems enable the process abstraction
- Linux API and its use. (Key idea: not implementation of API)

Memory Hierarchy



Linking

 The process of collecting and combining various pieces of code and data into a single file that can be loaded (copied) into memory executed

Topics

- Static and dynamic linking
- Object files, relocatable code
- Symbols, symbol resolution, symbol tables
- Position independent code
- Library interpositioning

Virtual Memory

 Illusion that a program has the entire physical address space for its own use even though many programs may be co-running

Topics

- Address translation
- Translation-lookaside buffers
- Page tables and page fault
- Dynamic storage allocation
- Garbage collection

System-Level I/O

 Managing storage device (e.g., disk) as a reliable and easy-touse persistent storage resource

Topics

- How to use the Linux filesystem API
 - Not a course for learning to implement filesystems
 - Appropriate API usage is an art in its own right!
- System call and memory-mapped I/O
- Includes aspects of virtual memory

Network Programming

- High-level and low-level I/O contd., with extension to network programming
- Very similar API for storage and networking I/O
- Internet services, web servers

Concurrent Programming

- Concurrent server design
- Threaded server versus process-based server
 - Last year's assignment's key theme
- I/O multiplexing with select
- Some aspects of parallel programming
 - Stepping-stone to parallel systems course

Java Virtual Machine (JVM)

- Java programming language has an entire runtime to deliver on its key promises
 - Memory safety + portability
 - Nothing comes for free in systems!
- We will cover fundamentals of JVM internals
- Will inform us why Java is slower than C and what can be done about that
 - Virtual machines is a powerful idea!



Big Data Frameworks

- Big data frameworks today process very large datasets
- They stress every aspect of key COMP2310 topics
 - Memory
 - Storage
 - CPU
 - Concurrency and networking
- We will study a selection of datacenter frameworks
 - Lucene search engine, RocksDB key-value store, Redis cache, Spark for machine learning analytics

Big Data Frameworks

 Typical data processing framework you can aim to implement after COMP2310

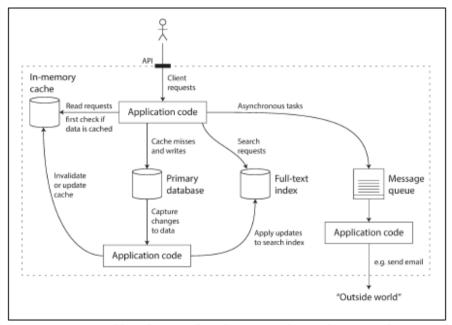


Figure 1-1. One possible architecture for a data system that combines several components.

From book: Designing Data Intensive Applications, Page 5

Assessment

Checkpoint 1

- Reverse engineering x86-64 object code
- Proficient in low-level C programming
 - pointers, string manipulation, etc

Assignment 1

- Implementing a memory allocator or malloc() from scratch
- Open-ended extensions on top of a base spec

Checkpoint 2

- Concurrency fundamentals
- Pthread synchronization

Assignment 2

- Related to networking with aspects of concurrency
- Last year assignment was a web proxy with a user-level cache
- Some changes this year but similar in inspiration

Quiz 1

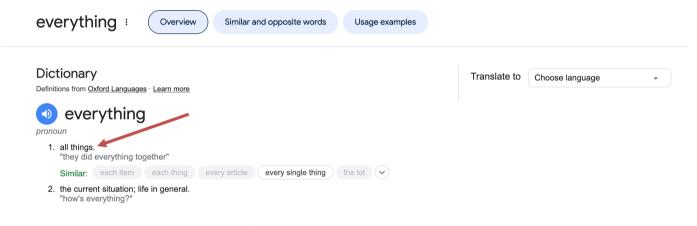
- Processes and signals
- Tests first four weeks of content
- Lab # 4 content is assessed indirectly by the first quiz

Quiz 2

- Memory allocation, virtual memory, cache, storage, some database concepts
- Tests weeks 4 7 content
- Lab # 6 content is assessed indirectly by the first quiz

Final Exam

Everything!



- Inclusive of week 12
- Every lab
- Every slide

Assessment Schedule

- 2-week window to attempt quizzes
- 8 9 days for checkpoints
- ~ 2 weeks for assignments

	<u>Release</u>	<u>Due</u>
Checkpoint 1	Aug 6	Aug 14
Quiz 1	Aug 14	Aug 28
Assignment 1	Aug 23	Sep 6
Quiz 2	Sep 18	Oct 2
Checkpoint 2	Sep 27	Oct 6
Assignment 2	Oct 10	Oct 25

Breakdown

- Checkpoint 1 (5%)
- Checkpoint 2 (5%)
- Quiz 1 (2.5%)
- Quiz 2 (2.5%)
- Assignment 1 (20%)
- Assignment 1 (20%)
- Final Exam (45%)

Admin & Logistics

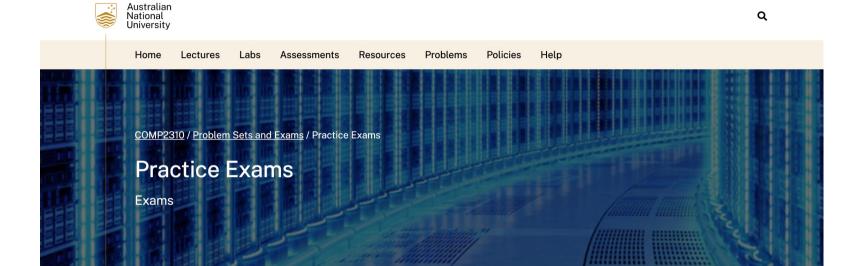
Succeeding in this course

- Pay attention to lecture content
- Finish all labs
- Read the textbook
- Submit all assessments

Assessment Difficulty

- Assignments are manageable if you start early
- Possibly the most "adventurous" exam of your ANU journey
- Check out the past year's exam and rubric on the website
- If you spend many hours finishing the first two labs and struggle with checkpoint 1
 - Make sure you finish COMP2300 first
 - Reconsider taking this course if it's not compulsory
 - Focus on the key points in the last slide

2022 Exam



Past Exams

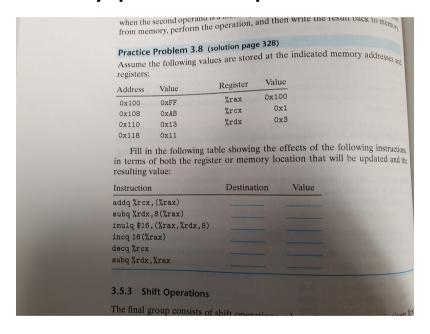
- 2022 Exam
- 2022 Exam (Solution/Rubric)

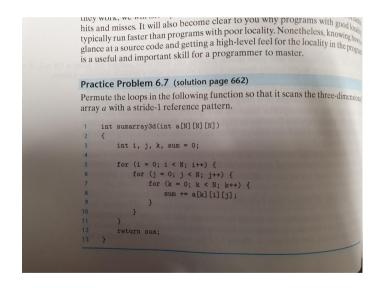
Readings: Book Chapters

Chapters	Topics/Weeks	2310 Coverage
1	COMP2300	Recommended
2	COMP2300	Not required
3	Weeks 1 – 2	Full except 3.11
4	COMP2300	Not Required
5	Weeks 1 – 2	Selected
6	Week 3	Full
7	Week 12	Full
8	Week 4	Full
9	Weeks 5 – 6	Full
10	Week 7	Full
11	Week 9 – 10	Full
12	Week 11	Full

Practice Problems

Try practice questions in book (answers in the book)





Cheating/Plagiarism

- Copying code, retyping by looking at a file
- Describing a solution to someone else so they can then type
- Searching the web for solutions to quiz or assignment
 - Last year's iteration of COMP2310, other universities' solutions in English or another language
- Copying from a github repository with minor or no modification
- Use of AI to generate your code

Cheating/Plagiarism

- Helping others by supplying code
- Debugging their code
- Telling them how to put together different code snippets to reach a working solution

Not Cheating

- Explaining how to use a tool
 - GDB, GCC, Valgrind, Editor, VSCode, Shell
- High level discussions
 - Not pseudo-code, not specific algorithms
- Using code supplied with the book
- Using Linux manpages
- Do not do this: COMP2310 malloc solution 2022

Cheating Consequences

- Action to uphold integrity begins at the time of discovery (not at the end of course)
- Last year, I read all submitted code from every student (but we will use automated tools as well)
- Some students were unable to pass the course due to academic integrity
- **Bottomline:** We want you to get the experience of dealing with systems programming issues from scratch!

Tutorials/Labs

- Labs are a critical component of this course (one every week)
- Handout will be posted on the website "Labs" before each lab
- First 2 labs
 - Becoming comfortable in C: pointers, bit-level manipulation, malloc() / free()
 - Lab 3 is assessed as the first checkpoint (no help from tutors)
- Lab 4
 - Process API and signal handling
- Lab 5 6
 - We will teach you to write a basic memory allocator, i.e., implementation of malloc()
- Lab 6 (assignment 1 week)
- Lab 7
 - Storage I/O
- Lab 8
 - Concurrency fundamentals
- Lab 9
 - Concurrency & Networking (sockets API)
- Labs 10, 11
 - Assignment 2, threads & concurrency (pthreads)

Networking &

Assignment Submission

- Extensions will be granted on a per-request basis
 - Via the extension app
- Assignment submissions are handled via Gitlab
 - You will learn more about it in the labs
 - Make a habit of using Git properly
 - Push often, always pull the latest

Each student submits their own work. No groups.

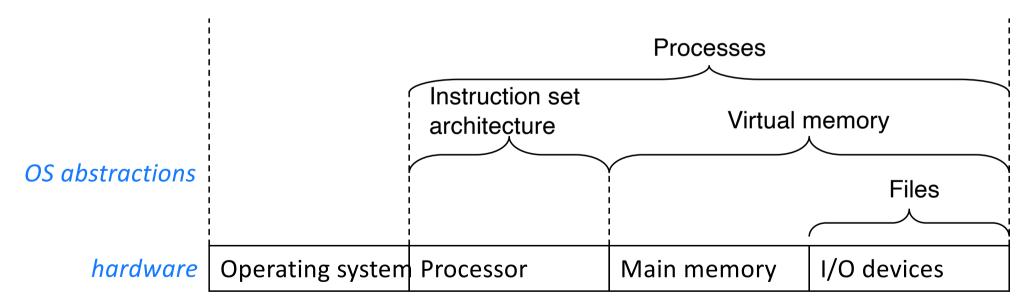
Note that: Student + AI = Group

Rough Plan for Lectures

- 1. Overview and x86 assembly
- 2. Optimizations and security implications (x86 as a vehicle)
- 3. Memory hierarchy
- 4. Processes and signals (abstraction for CPU, memory, I/O)
- 5. Virtual Memory (abstraction for main memory)
- 6. Dynamic memory allocation (memory allocator design)
- 7. Big data frameworks that are memory and I/O intensive
- 8. Storage and File I/O (abstraction for I/O devices)
- 9. Networking
- 10. Concurrency
- 11. Linking
- 12. Revision (time permitting!)

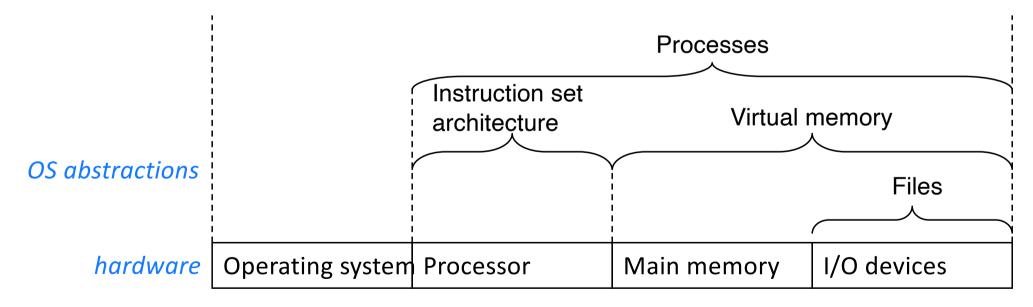
Course Organization (1)

- First 6 weeks lay the foundation of systems programming
 - They deal with CPU and memory virtualization
 - CPU and memory as a raw resource is not safe for multiuser systems and real programs



Course Organization (2)

- Next week deals with abstraction for storage I/O devices
 - Without storage and files, no serious application can work
- Next week: puts everything together to discuss real-life big data processing frameworks



Course Organization (3)

- And finally, every system must communicate with other systems (world wide web)
 - We move to networking
 - Networking is also I/O so an extension of storage I/O
- A networked application must deal with multiple producers and consumers of information
 - In comes concurrency!
- Finally, we end the course with linking (how large programs that use external libraries are compiled efficiently and safely)
 - Ideally fits in week 4, but we need to approach memory early

Welcome and Have Fun!