

Basic Blocks and Traces

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Basic blocks and Traces

Issues:

- To simplify translation there are mismatches between tree code and actual machine instructions:
 1. CJUMP to two labels; machine conditionals fall through on false
 2. ESEQ and CALL order evaluation of subtrees for side-effects – constrains optimization
 3. CALL as argument to another CALL causes interference between register arguments
- Can rewrite equivalent trees without these cases:
 - SEQ can only be subtree of another SEQ
 - SEQs clustered at top of tree
 - might as well turn into simple linear list of statements
- 3-stage transformation:
 1. to linear list of *canonical trees* without SEQ/ESEQ
 2. to *basic blocks* with no internal jumps or labels
 3. to *traces* with every CJUMP immediately followed by false target

Canonical trees

1. No SEQ or ESEQ
2. CALL can only be subtree of EXP(...) or MOVE(TEMP t,...)

Transformations:

- lift ESEQs up tree until they can become SEQs
- turn SEQs into linear list

$$\begin{array}{lcl} \text{ESEQ}(s_1, \text{ESEQ}(s_2, e)) & = & \text{ESEQ}(\text{SEQ}(s_1, s_2), e) \\ \text{BINOP}(op, \text{ESEQ}(s, e_1), e_2) & = & \text{ESEQ}(s, \text{BINOP}(op, e_1, e_2)) \\ \\ \text{MEM}(\text{ESEQ}(s, e_1)) & = & \text{ESEQ}(s, \text{MEM}(e_1)) \\ \text{JUMP}(\text{ESEQ}(s, e_1)) & = & \text{SEQ}(s, \text{JUMP}(e_1)) \\ \text{CJUMP}(op, \text{ESEQ}(s, e_1), e_2, l_1, l_2) & = & \text{SEQ}(s, \text{CJUMP}(op, e_1, e_2, l_1, l_2)) \\ \text{BINOP}(op, e_1, \text{ESEQ}(s, e_2)) & = & \text{ESEQ}(\text{MOVE}(\text{TEMP } t, e_1), \\ & & \text{ESEQ}(s, \text{BINOP}(op, \text{TEMP } t, e_2))) \\ \text{CJUMP}(op, e_1, \text{ESEQ}(s, e_2), l_1, l_2) & = & \text{SEQ}(\text{MOVE}(\text{TEMP } t, e_1), \\ & & \text{SEQ}(s, \text{CJUMP}(op, \text{TEMP } t, e_2, l_1, l_2))) \\ \text{MOVE}(\text{ESEQ}(s, e_1), e_2) & = & \text{SEQ}(s, \text{MOVE}(e_1, e_2)) \\ \text{CALL}(f, a) & = & \text{ESEQ}(\text{MOVE}(\text{TEMP } t, \text{CALL}(f, a)), \\ & & \text{TEMP}(t)) \end{array}$$

Taming conditional branches

1. Form *basic blocks*: sequence of statements always entered at the beginning and exited at the end:
 - first statement is a LABEL
 - last statement is a JUMP or CJUMP
 - contains no other LABELs, JUMPS or CJUMPs
2. Order blocks into *trace*:
 - every CJUMP followed by false target
 - JUMPs followed by target, if possible, to eliminate JUMP

Basic blocks

Control flow analysis discovers basic blocks and control flow between them:

1. scan from beginning to end:
 - LABEL l starts a new block and previous block ends (append JUMP l if necessary)
 - JUMP or CJUMP ends a block and starts next block (prepend new LABEL if necessary)
2. prepend new LABELs to blocks with non-LABEL at beginning
3. append JUMP(NAME done) to last block

Traces

1. Pick an untraced block, the start of some trace
2. Follow a *possible* execution path, choosing false targets first
3. Repeat until all blocks are traced

Cleaning up:

- CJUMP followed by true target: switch targets, negate condition
- CJUMP(o, a, b, l_t, l_f) followed by neither l_t nor l_f :
 1. create new l'_f
 2. rewrite as CJUMP(o, a, b, l_t, l'_f), LABEL l'_f , JUMP l_f
- JUMP l , LABEL $l \rightarrow$ LABEL l