

Hash functions \& Choosing a good hash futction

ANb-School of Computing - Structured Programming $1110 / 1140 / 6710$

## Hash Functions

A hash function is a function $f$ that maps a key $k$, to a value $f(k)$, within a prescribed range. It maps arbitrary sized keys to fixed-sized hashes.

A hash is deterministic. (for a given key, $k, f(k)$ will always be the same)


## Choosing a Good Hash Function

A good hash for a given population, $P$, of keys, $k \in P$, will distribute $f(k)$ evenly within the prescribed range for the hash.

A perfect hash will give a unique $f(k)$ for each $k \in P$.
(Perfect hash is rarely possible:
Pigeon hole principle.)

## Why value determinism and even distribution?

- Lets reword how we stated determinism a bit:
- Given $x, y$, if $x==y$, then $h(x)==h(y)$.
- It follows that (by contraposition):
- If $h(x)$ != $h(y)$, then $x$ != $y$
- Even though we cannot give positive result ( $x$ is $y$ ) confidently,
- We can for the negative result ( $x$ is not $y$ )


## Why value determinism and even distribution?

- Now lets suppose $h(x)$ gives an integer in range [0, 9]
- And suppose input is uniformly random
- With 10 values (or buckets), given inputs $x$ and $y$, we have $90 \%$ chance of deciding $x$ != $y$ in $O(1)$
- There is still a $10 \%$ chance of collision, but we have cut down our average workload of later stage by 90\%
- HashSet vs ArrayList
- More applications in C05


## Why so many different hashes?

- We outlined the basic properties we look for in a hash
- Deterministic
- This is fundamental, and by definition of a mathematical function
- No exception to this requirement
- Even/uniform distribution of output
- This is not as indisputable - we don't know what the distribution of input is like
- But we try to obtain this by guessing what the "usual" input looks like, e.g. statistical analysis of past usage
- The second point is roughly where the divergence begins


## Why so many different hashes?

- For each input distribution, we would need a different hash function to get an even distribution

Evenly distributed output if input is normally distributed<br>Deterministic

| Evenly distributed |
| :---: |
| output if input is |
| evenly distributed |
| Deterministic |

Evenly distributed output if input is bimodal
Deterministic

## Why so many different hashes?

- Even more variations if we want additional properties

|  | Secure <br> Low memory usage |  |
| :---: | :---: | :---: |
|  |  |  |
|  | Evenly distributed output |  |
| Evenly distributed output if input is normally distributed | Deterministic | Very high memory usage |
| Deterministic | Low memory usage | Very slow |
|  | Evenly distributed output if input is evenly distributed | Evenly distributed output |
|  |  | Deterministic |
| Functions | Deterministic |  |

## Assume whatever distribution, pick a recipe

- From "Effective Java", Josh Bloch
- (An approximate translation below in pseudo code)
- Assume you have fields (or more generally values) field0, field1, field2, ...
- int result = 0; // accumulator
for (var field : fields) \{
var $x$ = convertToInt(field); // recursively call this hash if needed result = 31 * result + x;
\}
- How does this work? Suppose we have fields: $x 0, x 1, x 2$
- $\quad$ After loop 0 , result $=x 0$
- After loop 1, result $=31$ * $x 0+x 1$
- After loop 2, result = 31 * (31 * x0 + x1) + x2 = 961 * x0 + 31 * x1 + x2


## Intuition behind this pattern

- Why 961 * x0 + 31 * x1 + x2 (or similar)
- Each factor is used to disperse the field to a different band/partition of the output range
- So it is sensitive to change of any field



## Why 31?

- From the book, multiplication with 31 is very efficient:
- 31 * $x=(x \ll 5)-1$
- A more impactful answer (my guess) is we don't use odd prime very often. Suppose we use 100 instead of 31 :
- 10000 * x0 + 100 * x1 + x2
- Suppose we reduce the range of hash by doing \% 10, above becomes
- x2


## Why 31?

- Of course if we modulo 31, then we run into the same problem
- But not a super common number to use
- We see a lot of things using base 10, e.g. 10, 100, 1000
- Natural to human
- Or base 2, e.g. 1024, 2048, 4096
- Natural to machine
- Odd primes, less so. (We could have replaced 31 with 7 etc.)


## Converting things into int

- Again mostly based on the recipe from Effective Java book
- Any numeric primitive type: multiply by prime, hashCode(), Float.floatToIntBits(x)
- Recursive: 31 * node.left.hashCode() + node.right.hashCode()
- Linear/array: treat each element as a field in previous recipe


## More complex hash

- We can always mix and match, and use the recipe as the base skeleton
- Suppose we parameterise the recipe as
- hash(int prime, List<int> fieldHashes)
- Examples:
- hash(31, fields in some order) // original reciple
- hash(31, fields in some order) + hash(7, fields in reverse order)

