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Finding Security Bugs in Java Programs using Datalog

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Program Agenda

- Java Security Issues
- Example: Caller Sensitive Methods
- Rapid Prototyping of Program Analyses in Datalog
- Security Analysis for Caller Sensitive Methods
- Experiments

Zero-day Vulnerability Market [1]

Large Ecosystem

- Market for buying exploits
- Multi-billion dollar industry
- “Write once, run anywhere”
- Java is platform independent

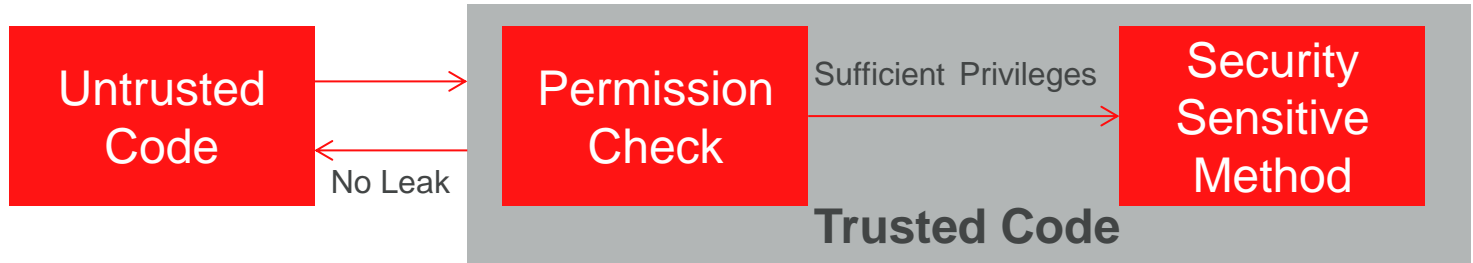
Software	Estimates in USD
Adobe Reader	\$5,000 - \$30,000
MAC OS X	\$20,000 - \$50,000
Flash or Java Browser Plug-Ins	\$30,000 - \$60,000
Microsoft Word	\$40,000 - \$100,000
Windows	\$50,000 - \$100,000
Firefox / Safari	\$60,000 - \$120,000
Chrome or Internet Explorer	\$80,000 - \$200,000
IOS	\$100,000 - \$250,000

[1] <http://www.net-security.org/secworld.php?id=12652> (March 2012)

Caller-Sensitive Methods (CSM)

One Possible Attack Vector for Java Exploits

- Security sensitive methods
 - must not be invoked unchecked on behalf of untrusted code
 - must not escape sensitive information
- If untrusted code invokes security sensitive methods
 - perform checks and prevent information leaks of sensitive information



Caller-Sensitive Methods (CSM)

Features and Issues

- 80% of JDK's public interfaces may directly or indirectly invoke a CSM
- Example of a CSM
 - `Class c = Class.forName("sun...")`
- CSM use reflection
 - hard to analyse
- Listed in Secure Coding Guidelines
 - Access Control / Section 9
- CSM use caller's class-loader or package access capabilities

Zero-day Exploit Example: CVE-2012-4681

Public method in `sun.awt.Toolkit`

```
public static Field getField(final Class klass,
                             final String fieldName) {
    return AccessController.doPrivileged(
        new PrivilegedAction<Field>() {
            public Field run() {
                try {
                    Field field = klass.getDeclaredField(fieldName);
                    field.setAccessible(true);
                    return field; ...
                }
            }
        }
    );
}
```

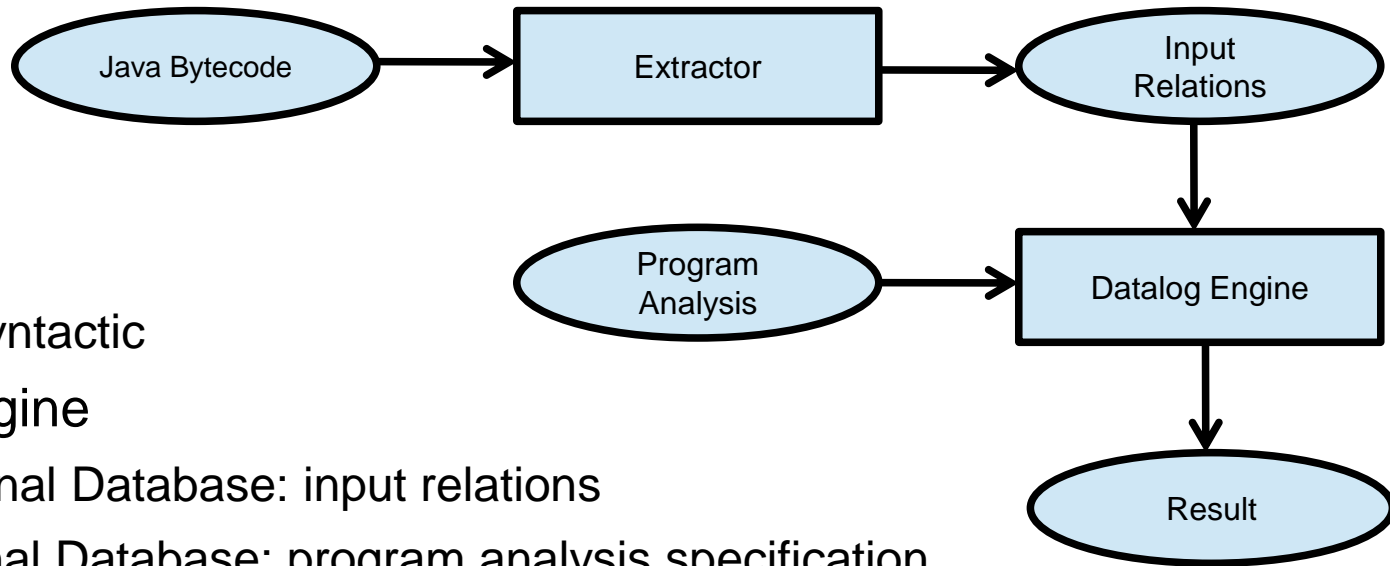

Finding Security Bugs

Automated Tools

- Testing for Security Bugs
 - Testing checks functional requirements and ***not security!***
 - Code-inspections are insufficient for finding (most) security bugs
 - Complex because of reflection, e.g., CSM
- Automated Bug-Checking Tools
 - Find security problems with static program analysis
 - By over-approximation using abstract interpretation
- Zero-day exploits ***demand rapid-prototyping capabilities***
 - Add new program analysis ***swiftly*** for new 0-day exploits

Bug Checker in Datalog

Framework



- Extractor
 - Purely syntactic
- Datalog Engine
 - Extensional Database: input relations
 - Intensional Database: program analysis specification

Security Analysis for CSM

Conditions

- Some conditions for causing security defects
 - Tainted inputs
 - User controls actual parameters of CSM
 - No permissions checks on a path from a public interface to CSM
 - Leak of sensitive information
- Building a security analyses in Datalog
 1. Points-to analysis
 2. Taint analysis (based on points-to analysis)
 3. All-path permission check
 4. Escape analysis (based on points-to analysis)

Points-To Analysis

Using Datalog

- Flow-insensitive, inclusion-based, context-insensitive (cf. J. Whalley'04)
- Abstract Domain
 - Variables
 - Local, actual/formal parameters, return-values, bases, this-variables
 - Heap-allocated objects
 - Creation-site as an abstraction for dynamically created objects
 - Heap-allocated object have fields
- Relations for computing points-to analysis
 - $vP(v,h)$: variable v may point to heap object h
 - $hP(h_1,f,h_2)$: field f of h_1 may point to h_2

Taint Analysis for CSM

Using Datalog

- Taint analysis tracks values emanating from tainted sources
 - Tainted values might be controlled by attacker
 - Tainted CSM parameters can be dangerous
- Taint analysis
 - Context- and flow-insensitive but object-sensitive
 - Public interfaces are a tainted source
 - Propagation rules for tainting objects
- Relations for computing taint analysis
 - $tH(h)$: heap object h might be tainted
 - $tV(v)$: variable v might be tainted

All-Path Permission Check

CSM

- On all paths from a public interface to a CSM callsite
 - A permission check must be performed (e.g. `checkPackageAccess`)
- CSM call-site could be exploited if,
 - a permission check is not performed on all paths, and
 - CSM parameters are tainted.
- Testing for all-path permission check
 - Classical dataflow analysis problem (e.g. GEN/KILL)
 - How to implement a dataflow analysis problem in Datalog?

All-Path Permission Check

Using Datalog

- The all-path permission check for a CSM call-site
 - **CheckedPaths(u)** $\Leftrightarrow \forall \pi \in \text{Path}(s, u): \exists v \in \pi: \text{Check}(v)$
 - **s** is a public interface
 - **u** is a statement (including CSM call-sites)
 - **Path(s,u)** is the set of all program path from **s** to **u**
 - **Check(v)** holds if statement **v** performs a permission check
- Dual logic because of Datalog semantics

$$\text{UncheckedPath}(u) \Leftrightarrow \exists \pi \in \text{Path}(s, u): \forall v \in \pi: \neg \text{Check}(v)$$

Experiments

Experiment

Problem Size

- OpenJDK 1.7
 - Number of variables: $\approx 1.5\text{M}$
 - Number of heap objects: $\approx 400\text{K}$
 - Number of methods: $\approx 170\text{K}$
 - Number of invocations: $\approx 600\text{K}$
 - Number of types: $\approx 18\text{K}$

Preliminary Results

Runtime & Effectiveness

Analyses	Time Taken
Basic (No Handling of Virtual Dispatch)	40 minutes
Virtual Dispatch + Call Graph Construction	7 hours

Intel i5-3320 (2.6GHz) machine with 16G memory running Ubuntu 12.10 using the BDDBDDDB engine

	Precision (%)	Recall (%)
First Taint Model	81	94
Second Taint Model	93	80

Using a reference implementation

Summary & Conclusion

Finding Security Bugs using Datalog

- Static program analysis is essential for checking security properties
- Implementation of program analysis using Datalog:
 - Rapid prototyping of different models
 - Extensible
 - Program analyses in Datalog are concise (=fewer bugs)
 - Debugging infrastructure still in its infancy
- Preliminary experiments
 - Datalog is efficient enough and effective

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