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# Points-To Analysis: Provenance Generation

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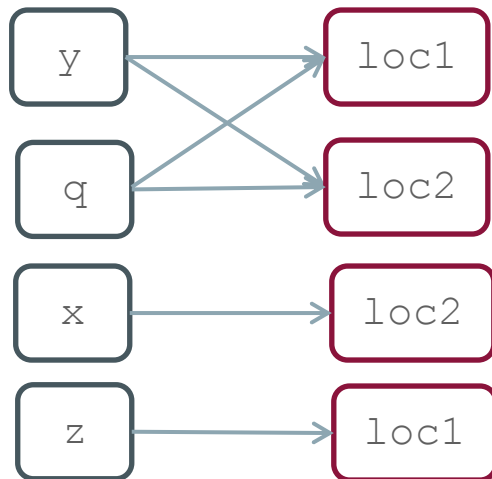
# Points-To Analysis

## Object-Oriented Languages

- Which variables “may point-to” which objects
  - Andersen: Flow-Insensitive, subset based
  - Steensgaard: Flow-Insensitive, equality based
  - Allocation site abstraction
  - Access-path abstraction
- Basis of other analyses
  - Taint, Escape
- [Tutorial by Yannis Smaragdakis](#)

# Andersen's Points-to Example

- Allocation site abstraction



- Scales for large code bases
  - JDK: 2M variables/500K allocation sites

```
public void bar() {  
    z = new T();    // loc1  
    foo(z);  
}
```

```
public void baz() {  
    x = new T();    // loc2  
    foo(x);  
}
```

```
private void foo (T q) {  
    T y = q;  
}
```

# Points-To Analysis

## Implementation

- Declarative specification: Datalog

```
VarPointsTo(x, obj) :- VarPoints(y, obj), Assign(x, y).
```

- Taint Analysis

```
Tainted(x) :- VarPointsTo(x, obj), TaintedObject(obj).
```

- Escape Analysis:

```
Escapes(obj) :- VarPointsTo(x, obj), ReturnVar(x, method),  
                Public(method).
```

- DOOP framework from Yannis Smaragdakis

– [Soufflé engine](#)

# Points-To Analysis

## Implementation

- Propagates/Uses points-to information
  - Reason/Provenance not explicit stored
    - ✓ `VarPointsTo(x, obj)`
    - ✗ `Assign(x, y), Assign(y, z), "z = new obj()"`
- Infeasible to explicitly store all provenance information
  - Preliminary Investigation: Each tuple requires about 50 more tuples for provenance

# Points-To Analysis

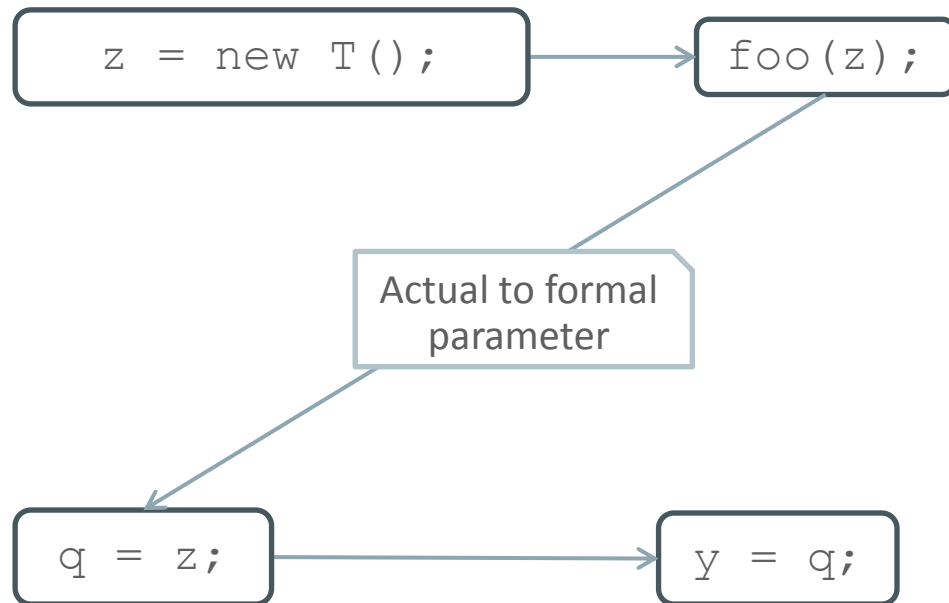
## Usability Issues/Challenges

- Debugging: Classify results as True Positives or False positives
  - Why does variable  $x$  point to object  $o$ ?
- Visualisation
  - Show me a taint/escape trace?
- Requires *provenance information*
- Challenge: Provenance generation for large code-bases
  - At least 1 million variables and 300 thousand allocation sites
  - Motivated by various subsets of the JDK



# Provenance

## Example



```
public void bar() {  
    z = new T();    // loc1  
    foo(z);  
}
```

```
public void baz() {  
    x = new T();    // loc2  
    foo(x);  
}
```

```
private void foo (T q) {  
    T y = q;  
}
```

# Problem Statement

- Given
  - Results of context sensitive points-to analysis
  - **Client** specified queries: Variable `var` points to location `loc`
- Find “*all provenance traces*” for each of the queries *using pre-computed results*
  - May points-to: All traces must be infeasible for report to be a False Positive
- Resource limits
  - Typical clients: Provenance generation less than 10% of points-to analysis

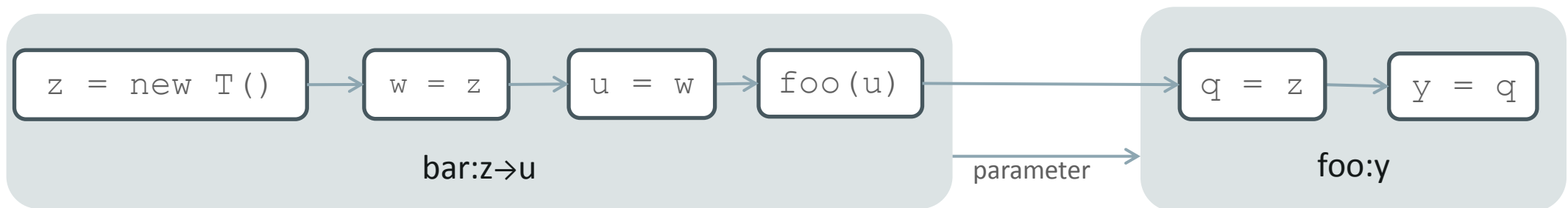
# Solution: Abstract Actual Trace

## Interprocedural Flow

- Track only flows between method boundaries

```
public void bar() {  
    z = new T();  
    w = z;  
    u = w;  
    foo(u);  
}
```

```
private void foo(T q)  
{  
    T y = q;  
}
```

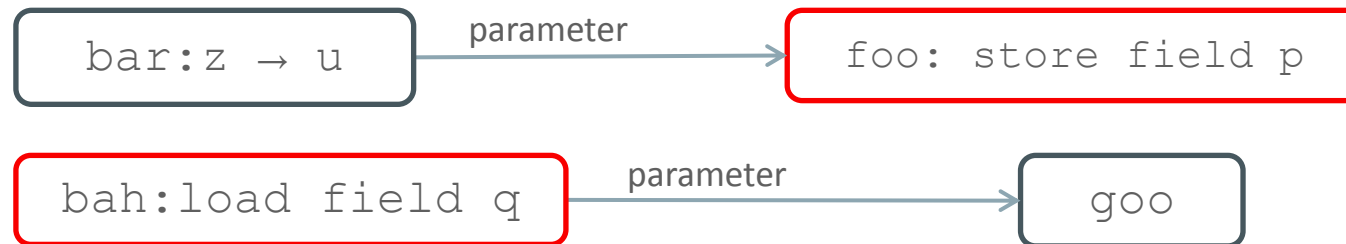


# Solution: Abstract Actual Trace

## Dataflow via Fields: Alias of base objects

- Load/Store also between method boundaries

```
public void bar() {  
    z = new T();  
    w = z;  
    u = w;  
    foo(u);  
}  
  
private void foo(T p) {  
    field = p;  
}  
  
public void bah() {  
    T q = field;  
    goo(q);  
}
```



# Dataflow via Fields: Alias

## Depth Limitation

- Alias of base variables can depend on alias of other variables
- Cascading alias
  - Expensive to compute
  - Hard to visualise
  - Increases FPs
- Solution:
  - Limit depth to 1 or 2
  - Sufficient for most traces

# Reusing Information

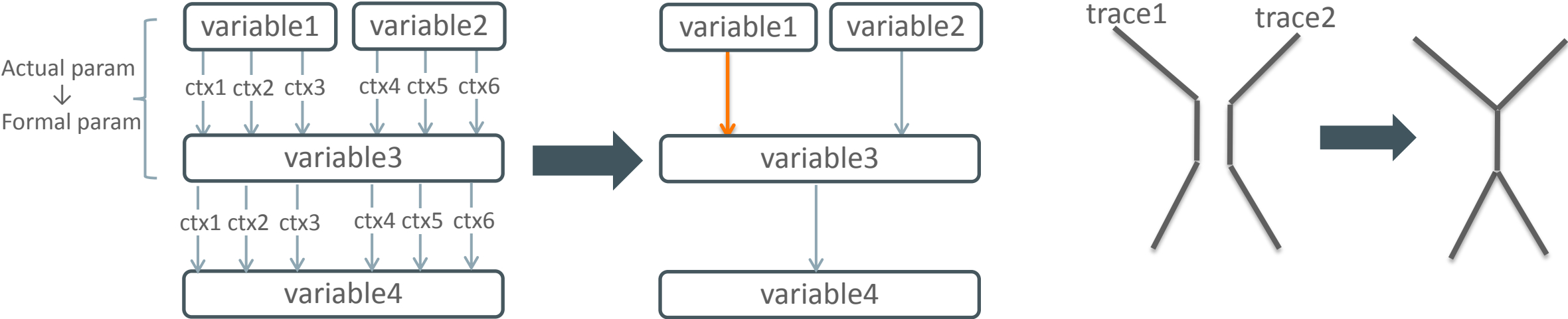
- Method boundaries: Identified using computed call-graph edge
- Load/Store pairs: Alias computed using points-to information
- Provenance: Identifies only such methods/variables/fields
- Optimisation: Remove irrelevant methods
  - Methods that return parameter

```
public T check(T v) {  
    if (v != null) return v;  
    throw new Exception("error");  
}
```

# Context Sensitivity

- Context in traces: Not scalable
- Solution: Contexts not part of the trace; but used to construct trace

$$\exists ctx, ctx', hctx: (hctx, loc, ctx, variable1) \in VP \ \& \ (hctx, loc, ctx', variable2) \in VP \ \& \text{CallGraphEdge}(ctx, \dots, ctx', \dots)$$



# Results

## Subsets of Various Versions of the JDK

Code-Base	VarPointsTo Size (tuples)	Client	#Client Queries	Provenance Graph			Time Overhead	Memory Overhead
				Nodes	Edges	Max Degree		
Subset Version-1	539Million	Taint	836	4287	8538	112	25%	3.4%
		Escape	10	13	14	2	1%	9.0%
Subset Version-2	871Million	Taint	900	4563	6774	105	23%	2.2%
		Escape	445	1115	6663	33	3.0%	8.8%



# Results: Alias Graph

## Context-Sensitive Points-To Analysis: Depth 1

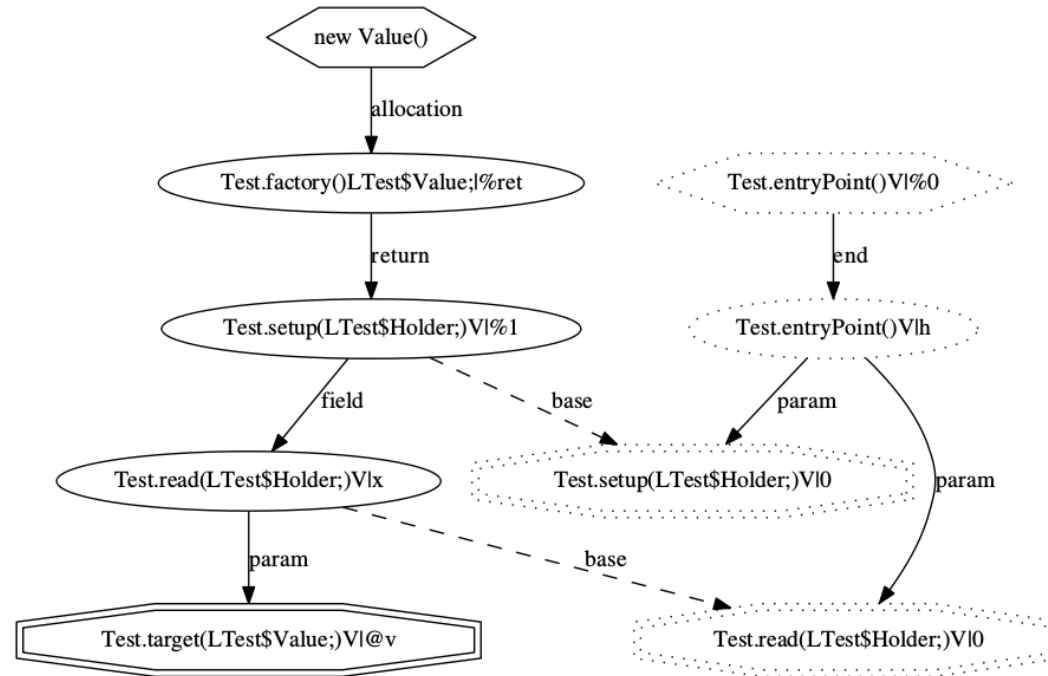
Code-Base	Client	Alias Graph		
		Nodes	Edges	Max Degree
Subset Version-1	Taint	23,791	98,245	136
	Escape	0	0	0
Subset Version-2	Taint	12,484	16,443	35
	Escape	7,588	28,180	37

# Presentation

- Graphviz
- Parfait: Internal format
- Shortest path
- Acyclic path
- Time/Output: Depends on size provenance graph
  - Taint analysis: Computing acyclic paths times-out

# Graphviz: Example

```
public class Test {  
    public static class Value {}  
    public static class Holder {  
        public Value v;  
    }  
  
    public static void entryPoint() {  
        Holder h = new Holder();  
        setup(h);  
        read(h);  
    }  
  
    private static Value factory() {  
        return new Value(); }  
  
    private static void setup(Holder h) {  
        h.v = factory(); }  
  
    private static void read(Holder h) {  
        Value x = forward(h.v);  
        target(x); }  
  
    private static void target(Value v) { ... }  
  
    private static Value forward(Value v) {  
        return v; }  
}
```



# Conclusion

- Provenance using sequence of method invocations
  - Limited control flow paths
- Reuse points-to for call-graph and alias analysis
- Scales for JDK
  - Time threshold of 10% not met: Further optimisations required
- Challenges
  - Support flow-sensitive analysis
  - Dynamic Traces: Querying mechanism of provenance information
  - Better UI for trace representation

# Questions

## Points-To Analysis: Provenance Generation

Contact [paddy.krishnan@oracle.com](mailto:paddy.krishnan@oracle.com) for more details

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